

# **Probabilistic Logic CNF for Reasoning**

**Song-Chun Zhu, Sinisa Todorovic, and Ales Leonardis**

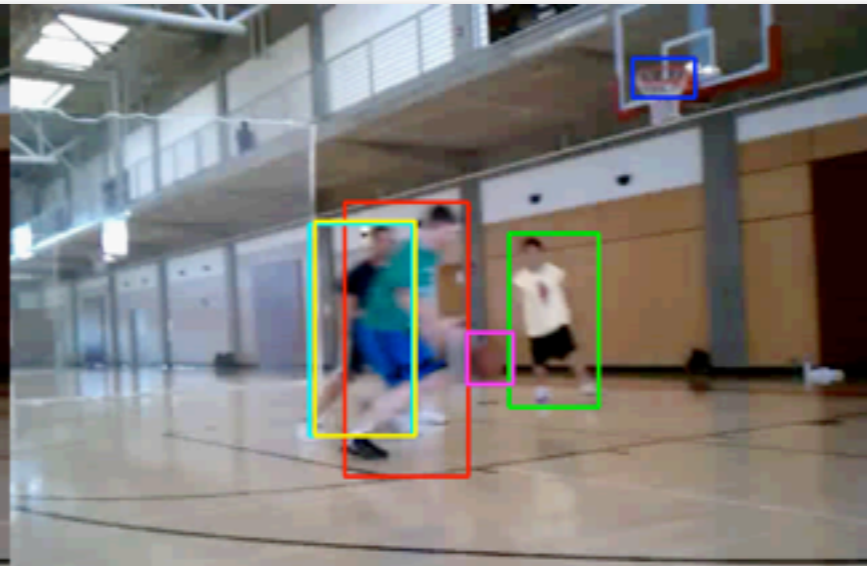
**At CVPR, Providence, Rhode Island  
June 16, 2012**

# Goal

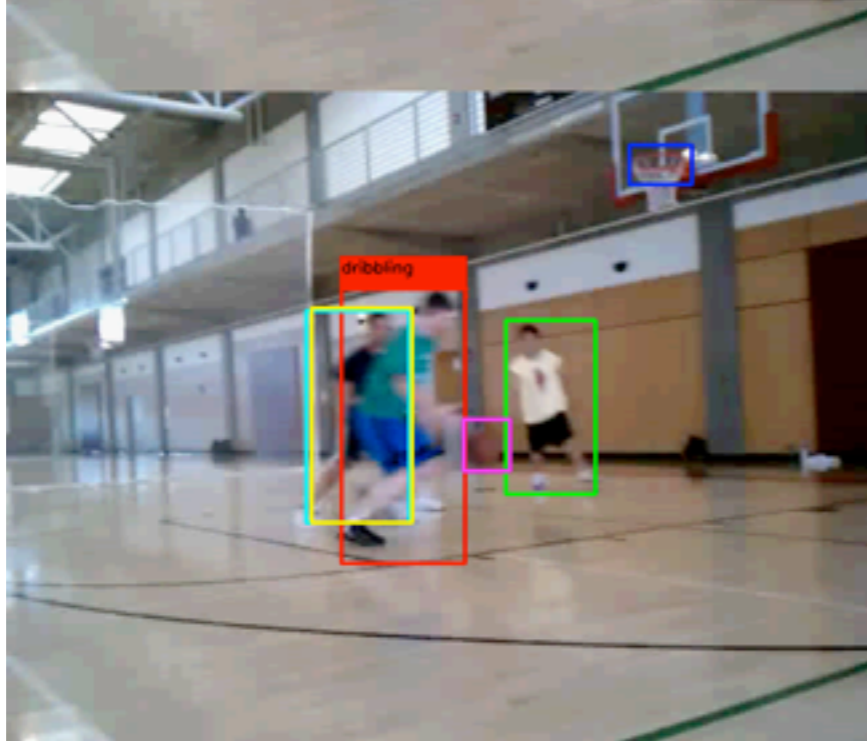
input



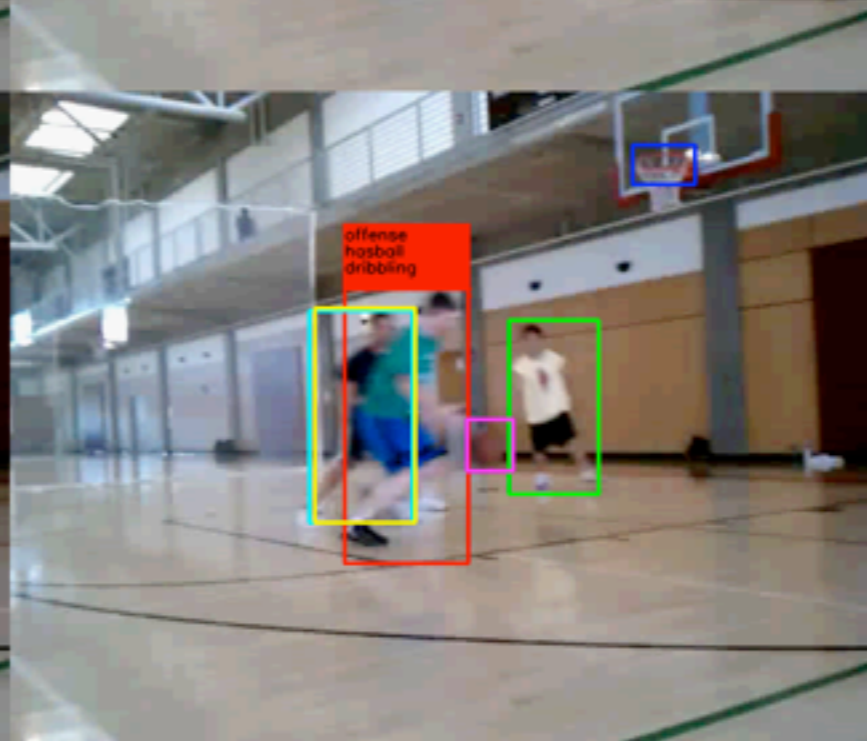
tracking



parsing

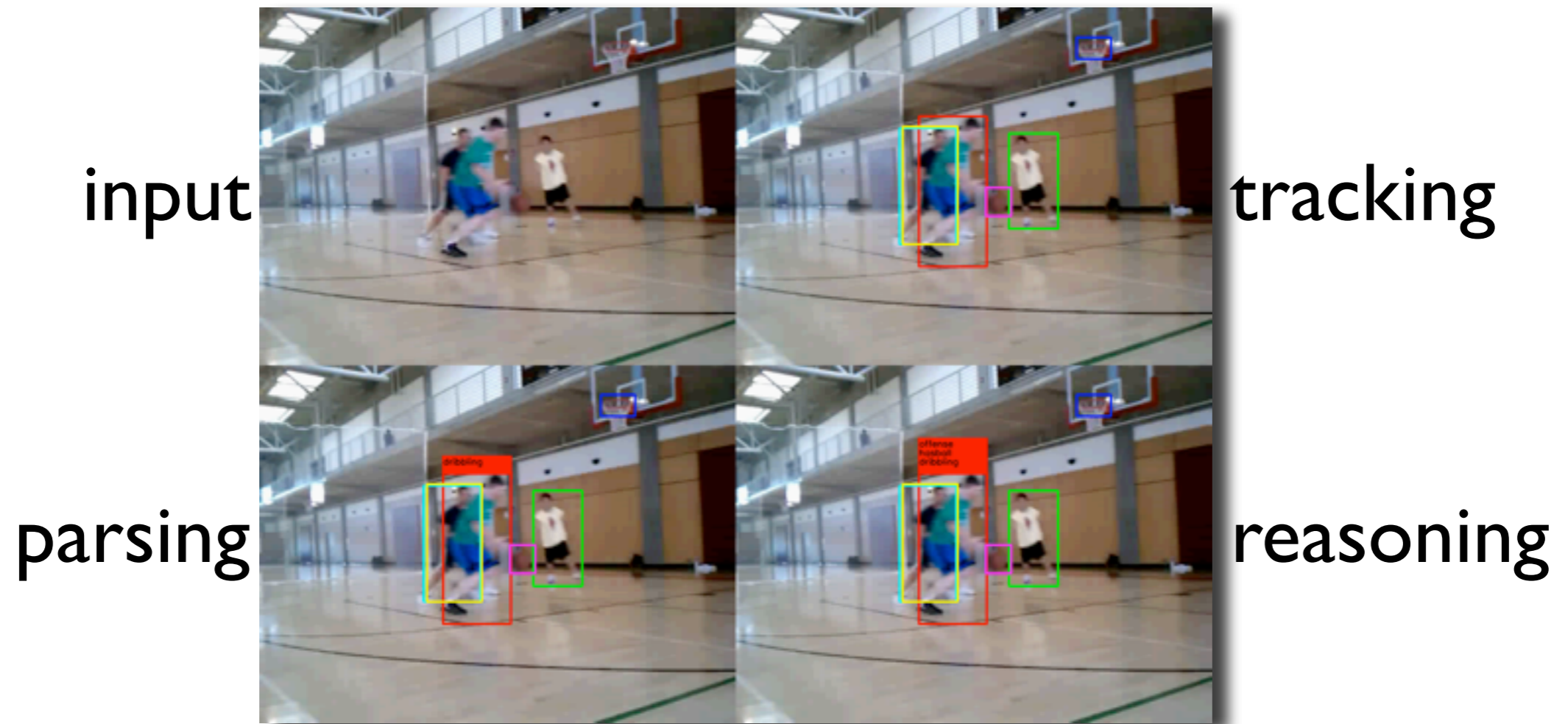


reasoning



activity recognition  
in structured domains

# Goal



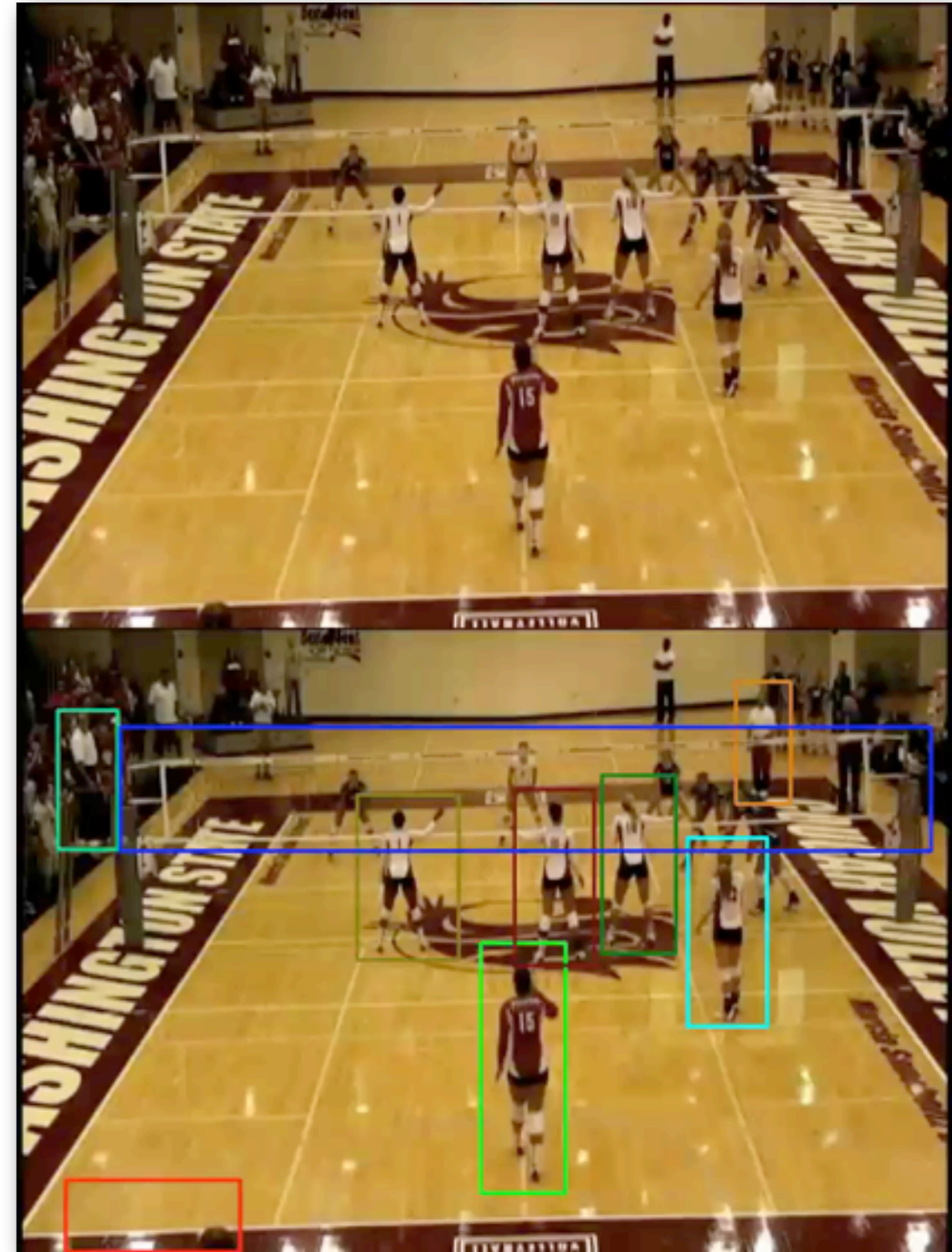
**what:** objects, events?

**where,** and **when?**

**why:** by explaining space-time relationships?

# Three Semantic Levels of Events

- Primitive actions:
  - single actor-object interaction
  - punctual actions
  - repetitive actions
- Activity:
  - Short-term human-human-object interactions
  - e.g., passing the ball, hugging
- Events:
  - Long-term interactions of a group of people and objects

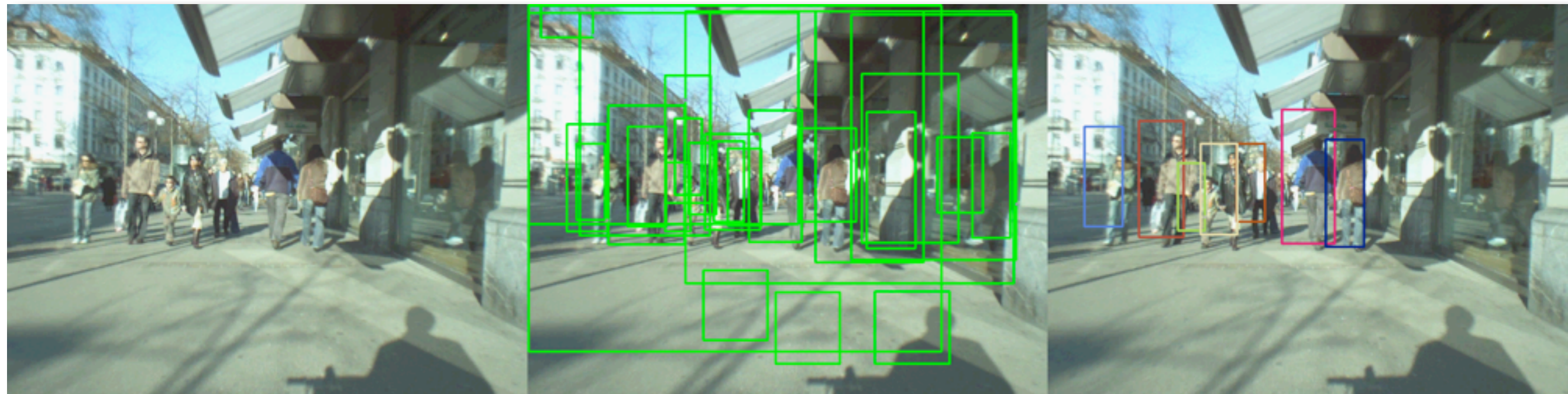


# In this Talk: Tracking & Parsing are Given

Input

Noisy detections

Our results



“Multiobject Tracking as Maximum Weight Independent Set”

CVPR 2011

# In Addition to Usual Challenges...

input



tracking

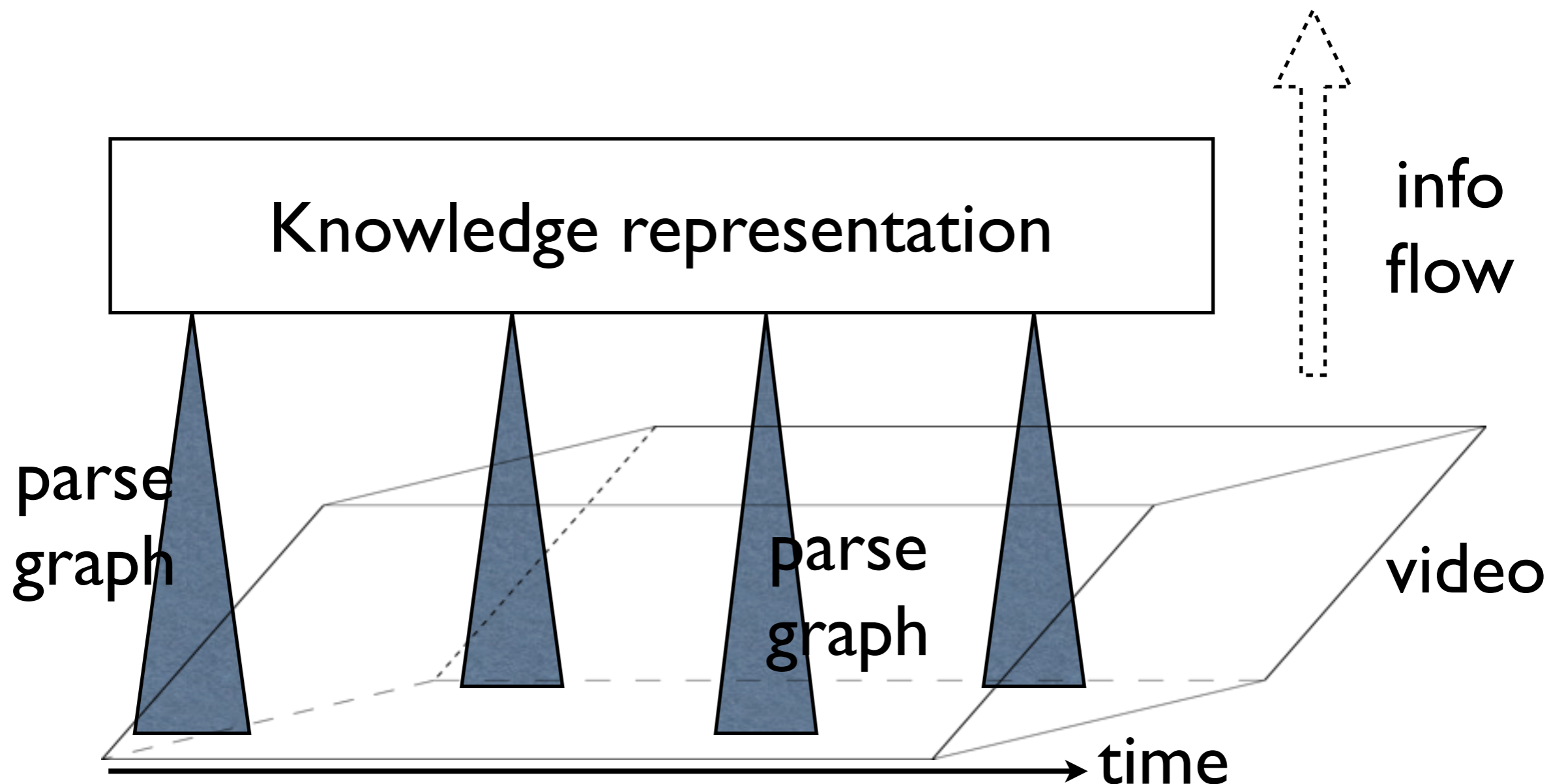
parsing

reasoning

- Who has the ball? -- partial occlusion
- Is the red player on offense? -- no direct cues
- Who violated the rules of basketball? -- domain

# Two Key Ideas

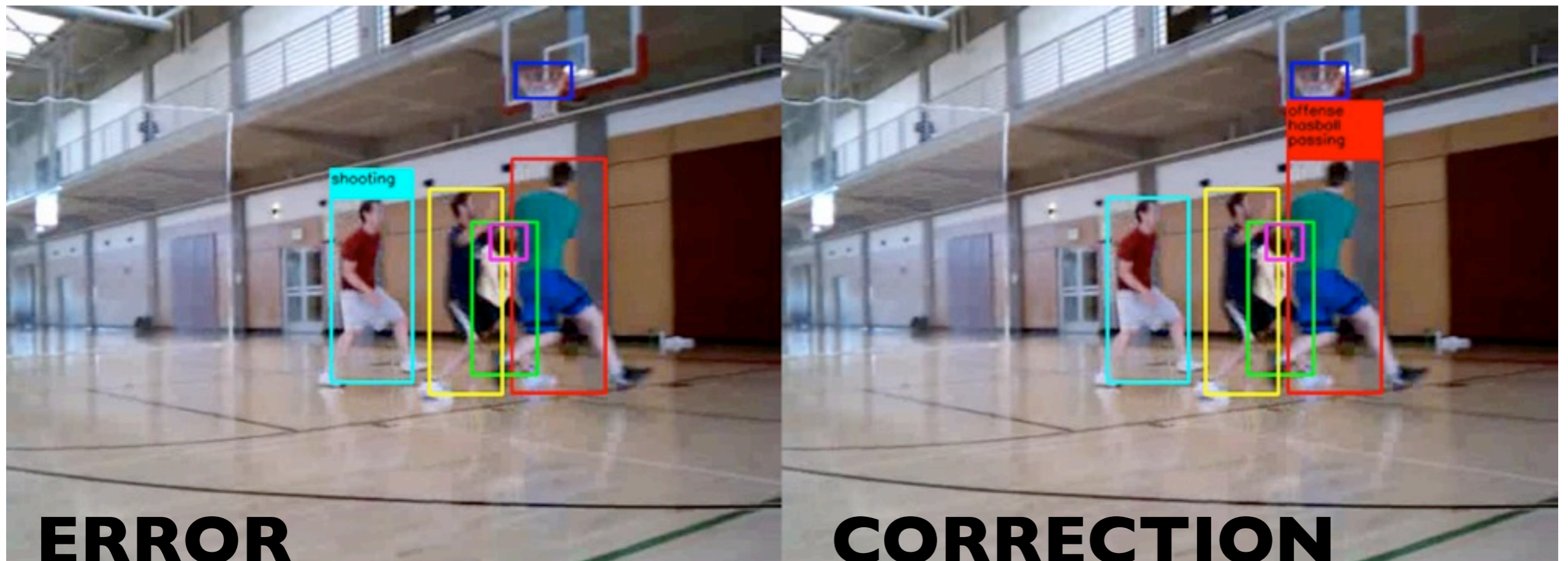
1. Ground reasoning onto parse graphs of primitive actions



# Two Key Ideas

## 2. Use domain knowledge to resolve uncertainty

Top-down correction of errors in tracking, parsing, entity resolution





# Knowledge Representation

- Probabilistic First-Order Logic
  - Rota & Thonnat '00 -- Declarative models
  - Siskind '01 -- Event logic
  - Nevatia et al. '04 -- Probabilistic ontology
  - Shet et al. '06 -- Multivalued logic
  - Richardson & Domingos '06 -- MLN
  - Shet et al. '07 -- Bilattice logic
  - Ryoo & Aggarwal '09 -- Space-time logic
  - Fern '09 - Penalty logic
  - Kersting & Raedt '11 -- Bayesian logic

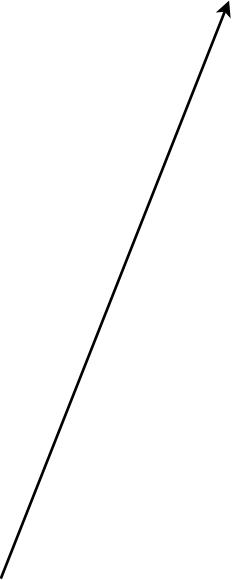
# Knowledge Base

$$\Sigma = \{(\phi_1, w_1), \dots, (\phi_n, w_n)\}$$

a set of weighted logic formulas

$$w_n = P(\phi_n @ I)$$

a distribution of costs  
of violating  $\phi_n$  over  
a time interval



# Logic Formula

$$\text{PassTo}(p, q) \rightarrow (\text{Pass}(p) \wedge_m \text{BallMoving} \wedge_m \text{Catch}(q))$$



Event symbol:  
e.g., interaction among a  
number of object types

# Syntax

$$\text{PassTo}(p, q) \rightarrow (\text{Pass}(p) \wedge_m \text{BallMoving} \wedge_m \text{Catch}(q))$$

Event symbol:  
e.g., interaction among a  
number of object types

Object types:  
e.g., person, tool

Short-hand notation for:

$$\neg\phi \vee \phi'$$

# Syntax

$$\text{PassTo}(p, q) \rightarrow (\text{Pass}(p) \wedge_m \text{BallMoving} \wedge_m \text{Catch}(q))$$

Event symbol:  
e.g., interaction among a  
number of object types

Three types of relations:

negation  $\neg\phi$

disjunction  $\phi \vee \phi'$


Temporal relations  
between time intervals  
where events are true

$$\phi \wedge_R \phi' \quad R \subseteq \mathbb{R}$$

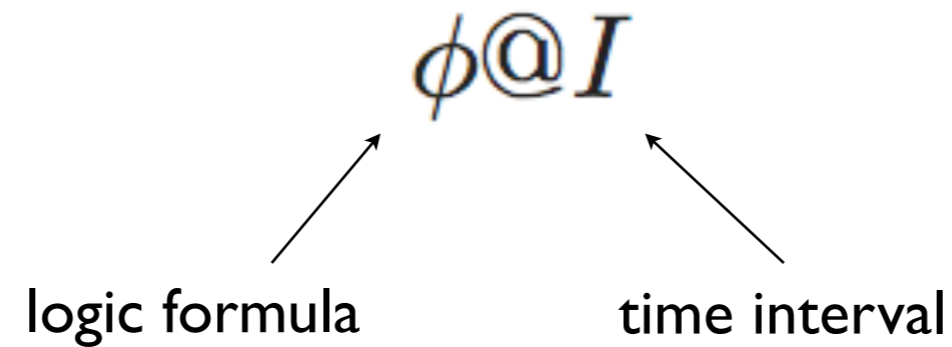
# Allen Temporal Relations

$I_1$	Relation	$I_2$	English	Definition	Inverse
$[m_1, m_2]$	<b>s</b>	$[n_1, n_2]$	starts	$m_1 = n_1$ and $m_2 < n_2$	<b>si</b>
$[m_1, m_2]$	<b>f</b>	$[n_1, n_2]$	finishes	$m_1 < n_1$ and $m_2 = n_2$	<b>fi</b>
$[m_1, m_2]$	<b>d</b>	$[n_1, n_2]$	during	$m_1 > n_1$ and $m_2 < n_2$	<b>di</b>
$[m_1, m_2]$	<b>b</b>	$[n_1, n_2]$	before	$m_2 < n_1$	<b>bi</b>
$[m_1, m_2]$	<b>m</b>	$[n_1, n_2]$	meets	$m_2 + 1 = n_1$	<b>mi</b>
$[m_1, m_2]$	<b>o</b>	$[n_1, n_2]$	overlaps	$m_1 < n_1 \leq m_2 < n_2$	<b>oi</b>
$[m_1, m_2]$	<b>=</b>	$[n_1, n_2]$	equals	$m_1 = n_1$ and $m_2 = n_2$	<b>=</b>

time interval



# Truth Values Assigned to Event Occurrences



observable  
event occurrences

$X$ : D-Dribbling( $P3$ )@[10, 30]

 from parse graphs

hidden  
event occurrences

$Y$ : Dribbling( $P3$ )@[20, 30]

 from reasoning

# Interpretation

$$(X, Y) \models (\text{HasBall}(P_1) \vee \text{HasBall}(P_2)) @ [10, 20]$$

an event occurrence is true

along interval [10,20]

in interpretation  $(X, Y)$



# Model

Given KB:  $\Sigma = \{(\phi_1, w_1), \dots, (\phi_n, w_n)\}$

$$P(X, Y | \Sigma) = \frac{1}{Z} \exp \left[ \sum_{\phi_n \in \Sigma} \sum_{I_i} P(\phi_n @ I_i) P_{pg}(\phi_n @ I_i) P((X, Y) \models \phi_n @ I_i) \right]$$

Prior distribution

confidence of the parse graph

likelihood that  $\phi_n$  is true in interpretation  $(X, Y)$

all time intervals where  $\phi_n$  is true

# Reasoning = Most Probable Explanation

$$(X^*, Y^*) = \text{MPE}(X, \Sigma) = \arg \max_{(X, Y)} P(X, Y | \Sigma)$$

We address intractable inference by:

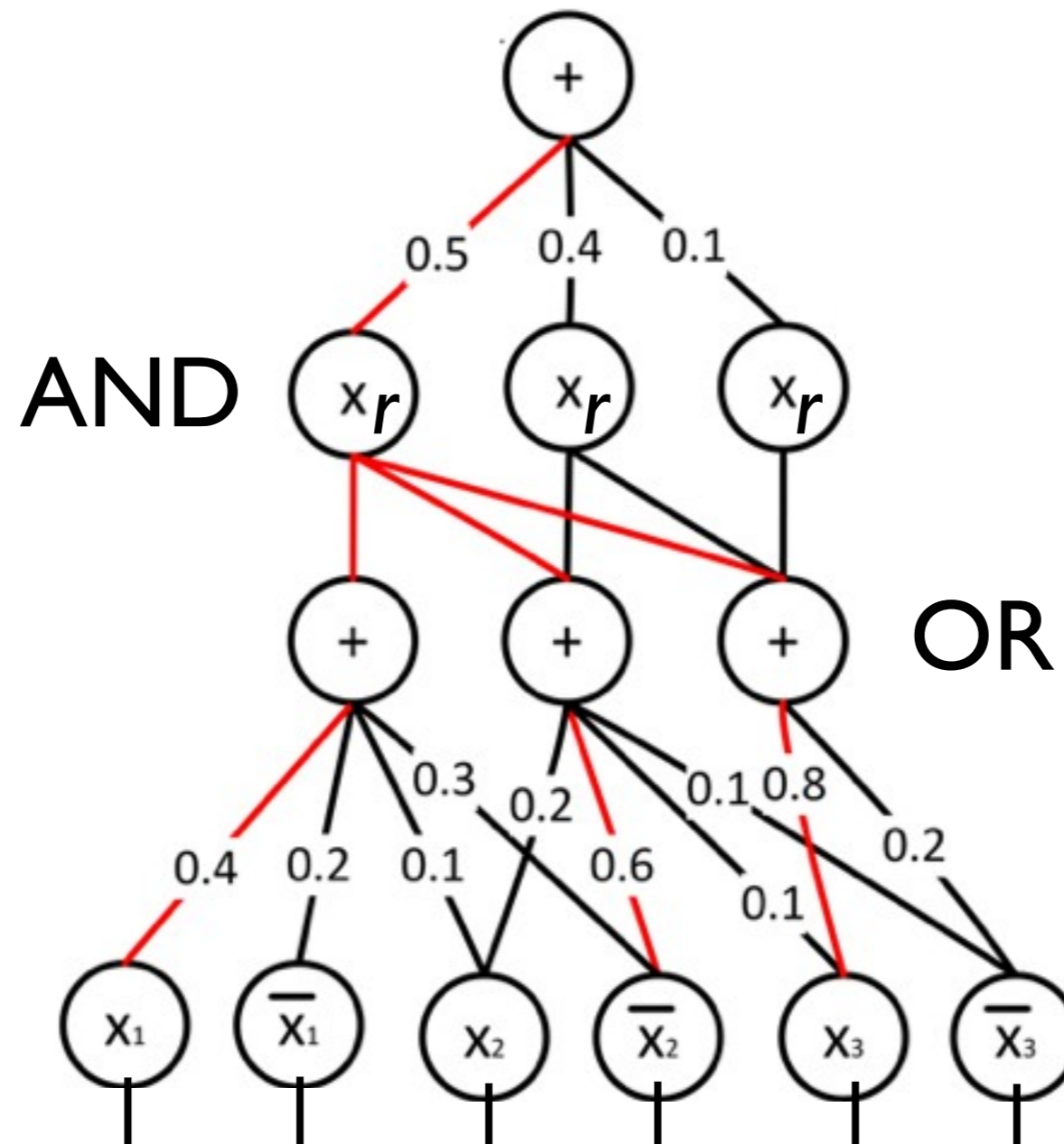
- Compiling  $\Sigma$  into CNF form  $\Rightarrow$  And-Or graph (AOG)
- Ensuring completeness and consistency of AOG
- Metropolis-Hastings moves over:
  - Logic formulas in  $\Sigma$
  - Arguments of the logic formulas
  - Time intervals along which the formulas are true

# Compilation of KB to AOG

Key idea:

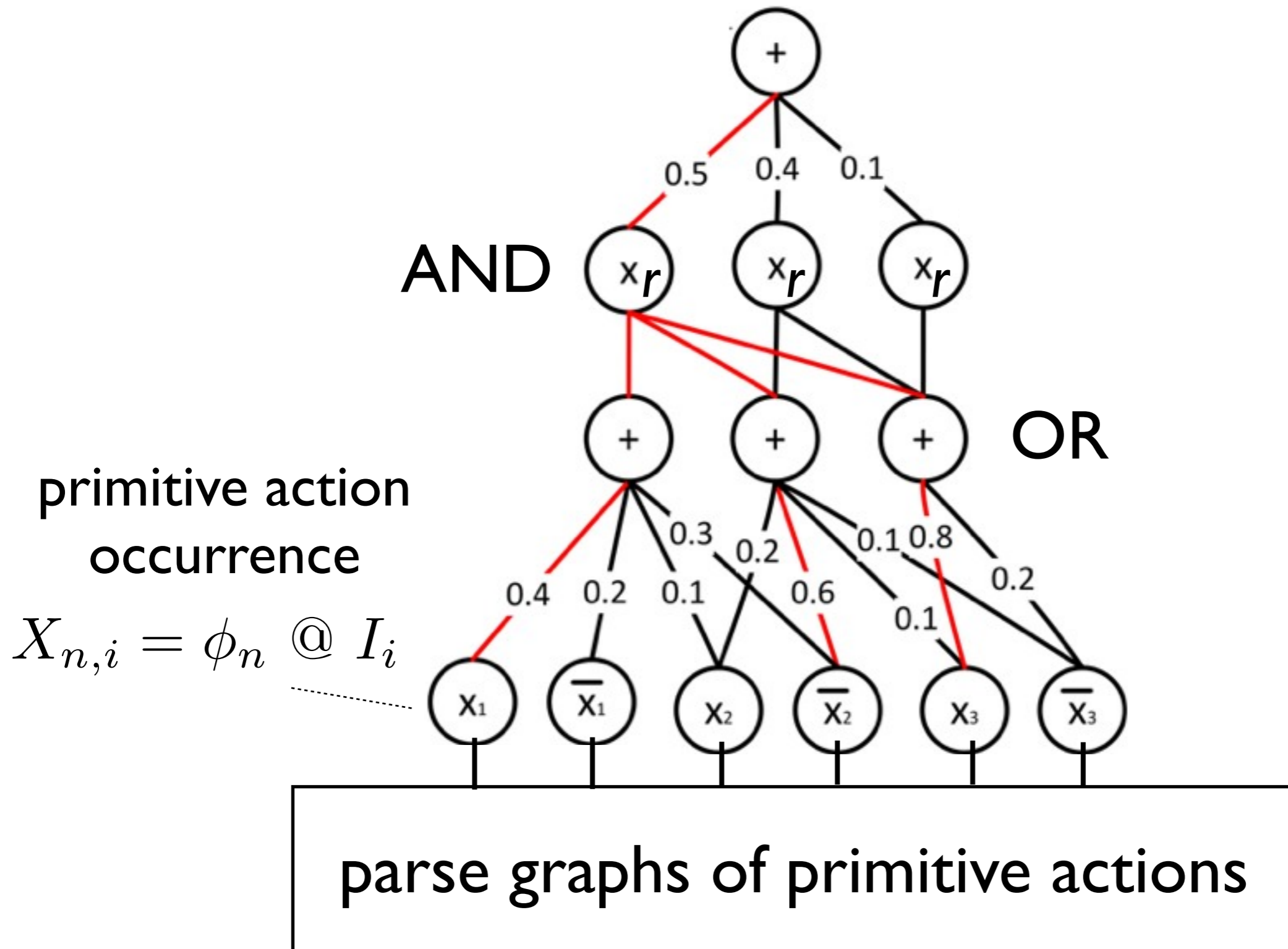
Arithmetic circuit -- Data structure for efficient inference  
Darwiche [2003]

# Compilation of KB to AOG

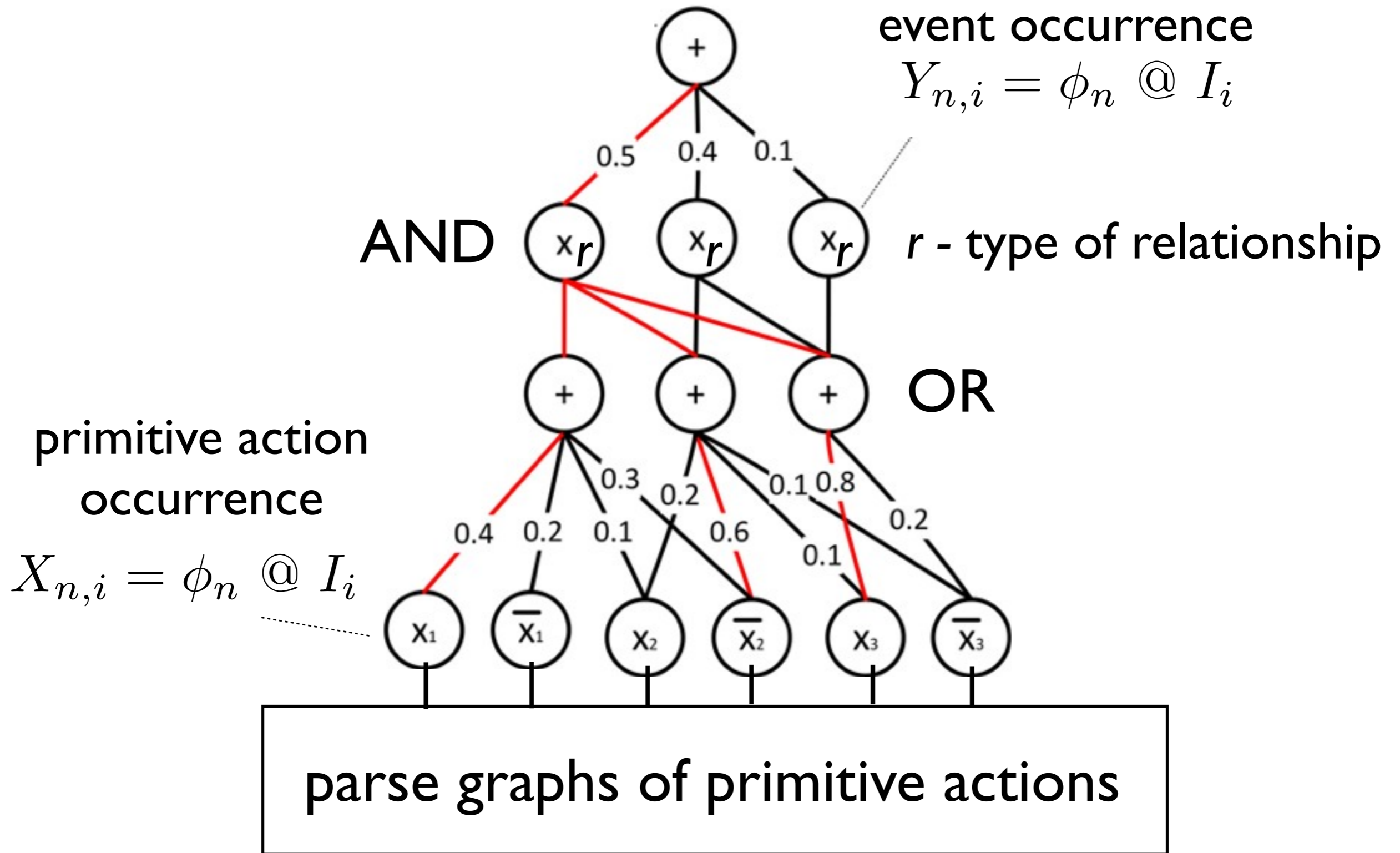


parse graphs of primitive actions

# Compilation of KB to AOG



# Compilation of KB to AOG



# Compilation of KB to AOG

$$w_n = P(\phi_n @ I)$$

event occurrence

$$Y_{n,i} = \phi_n @ I_i$$

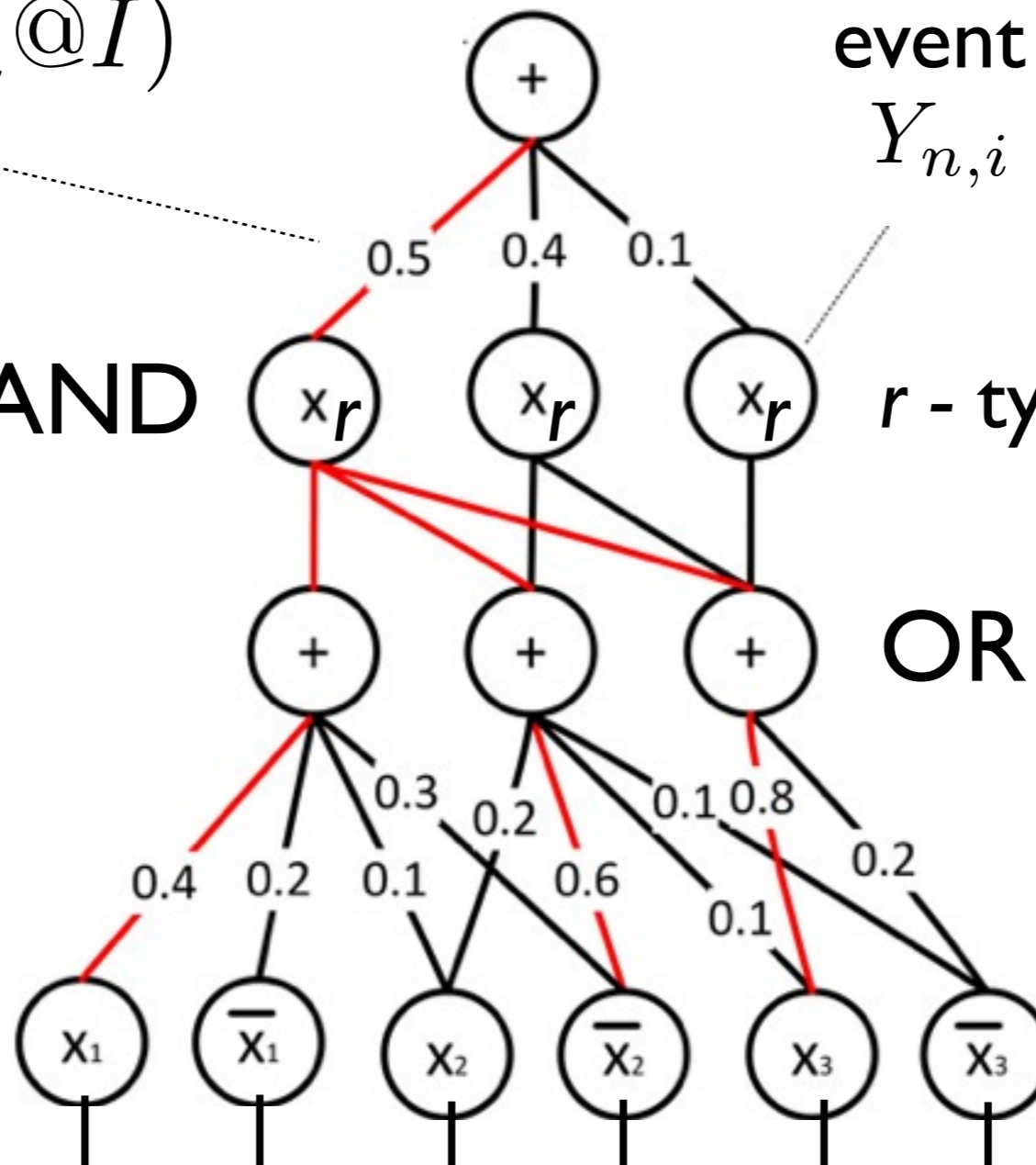
AND

$r$  - type of relationship

primitive action  
occurrence

$$X_{n,i} = \phi_n @ I_i$$

OR

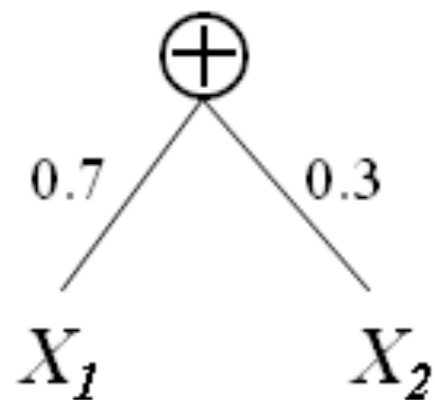


parse graphs of primitive actions

# Valid Compilation

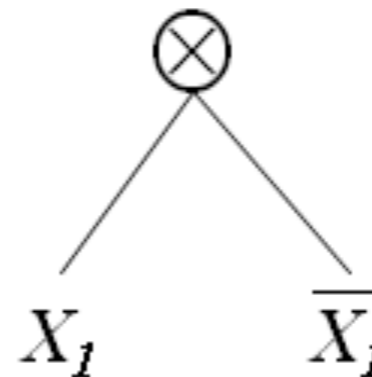
**Theorem:** AOG is valid iff it is **complete** and **consistent**

**Complete:** Under sum, children cover the same set of variables



**Incomplete**

**Consistent:** Under product, no variable in one child and negation in another



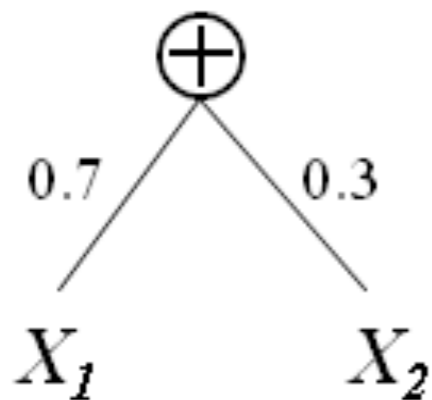
**Inconsistent**



# Efficiency

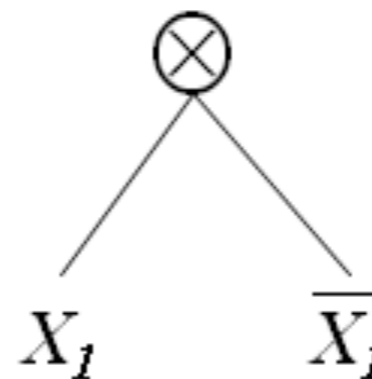
**Theorem:** Valid AOG allows polynomial inference in the number of nodes

**Complete:** Under sum, children cover the same set of variables



**Incomplete**

**Consistent:** Under product, no variable in one child and negation in another

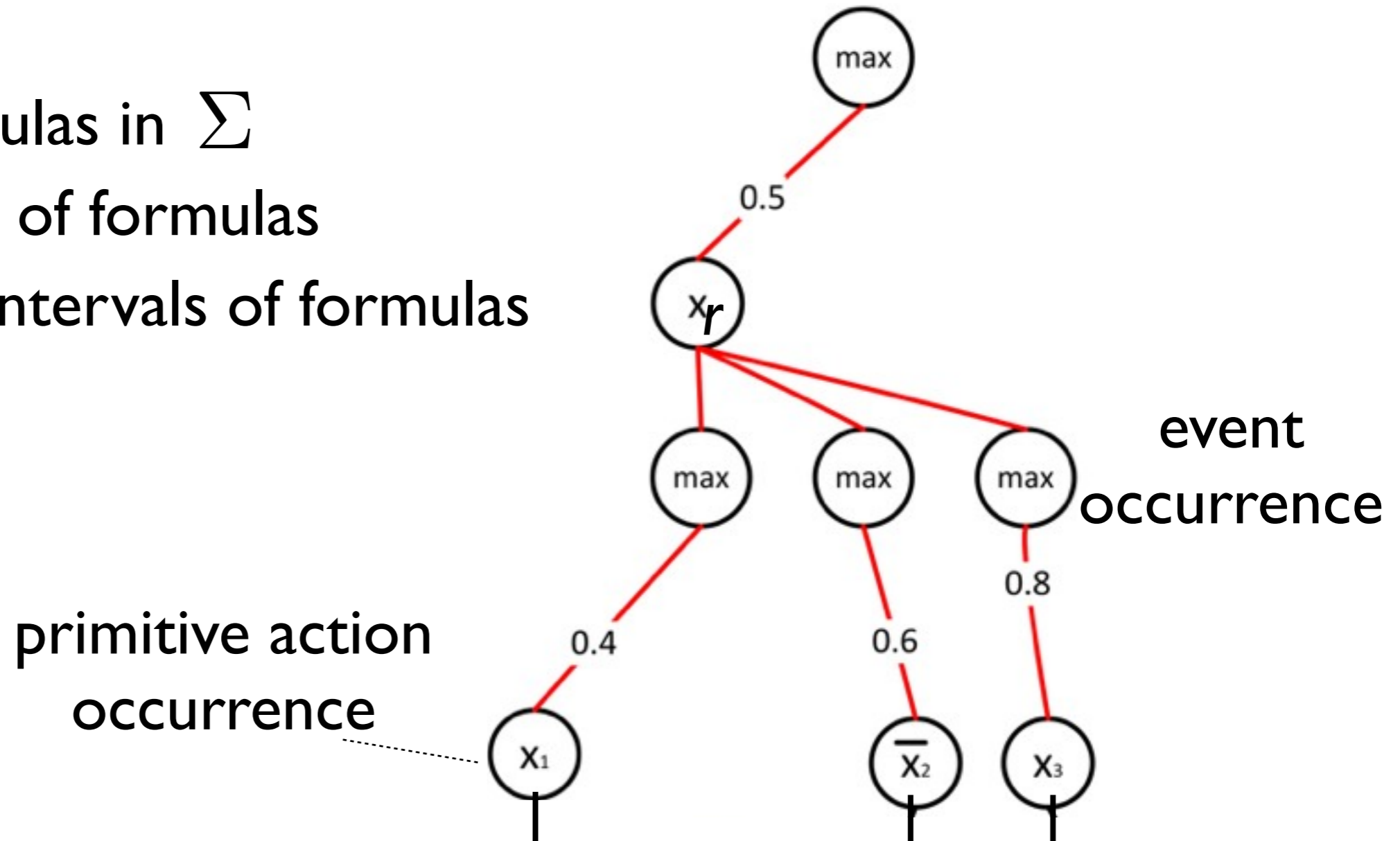


**Inconsistent**

# Most Probable Explanation

Identifies:

- Logic formulas in  $\Sigma$
- Arguments of formulas
- True time intervals of formulas



parse graphs of primitive actions

# Metropolis-Hastings Moves

two probable  
interpretations

$$A = (X, Y)_A \quad B = (X, Y)_B$$

$$\alpha(A \rightarrow B) = \min \left( 1, \frac{Q(B \rightarrow A)P(B|G)}{Q(A \rightarrow B)P(A|G)} \right)$$

proposal distribution

efficient proposals of time intervals  
without enumerating exponentially many  
subintervals of all intervals

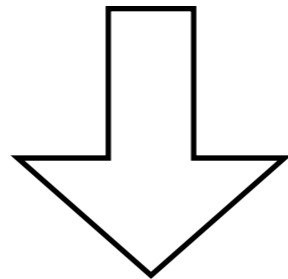
compiled KB  
into AOG

CVPR 2011

# Scheduling the Moves -- Open Problem

How to prioritize particular moves over:

- Logic formulas in  $\Sigma$
- Arguments of formulas
- True time intervals of formulas



Our approach:

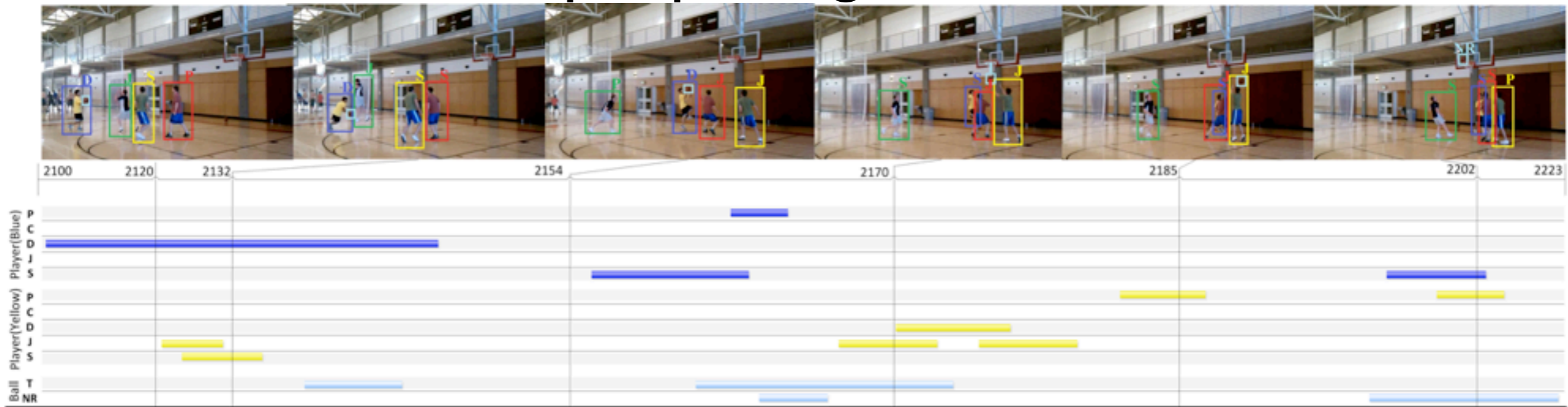
1. Map the current interpretation into a feature vector

$$(X, Y)_A \rightarrow \Psi_A$$

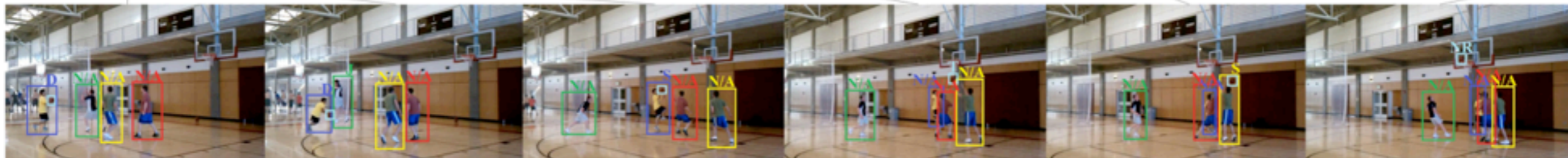
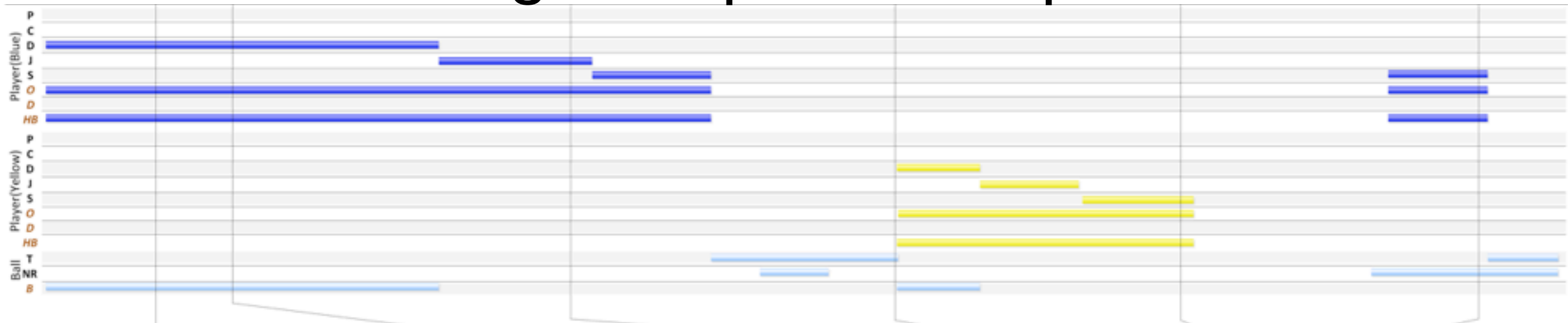
2. Classify the feature vector

# Results -- CVPR 11

input parsing results

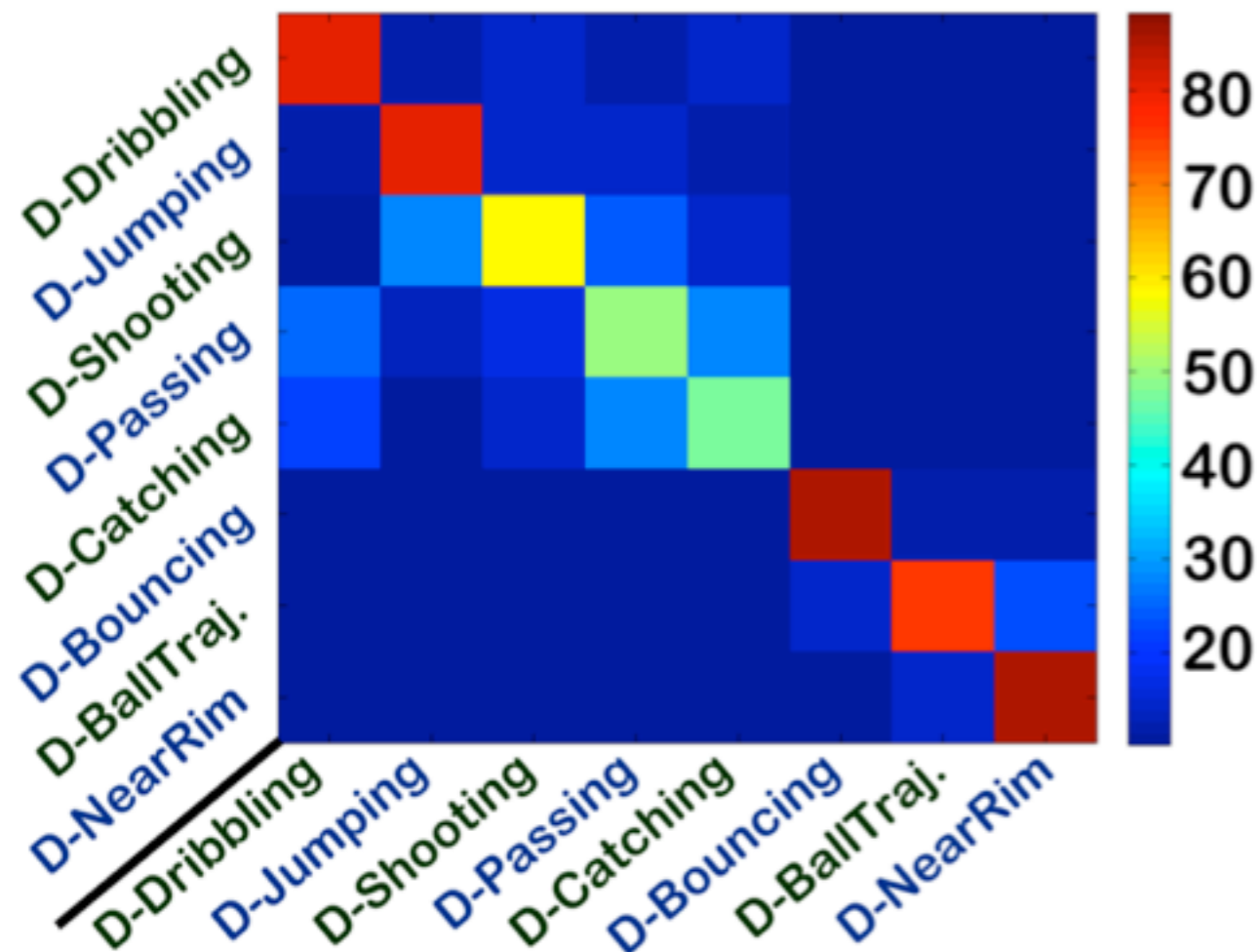


reasoning: most probable explanation

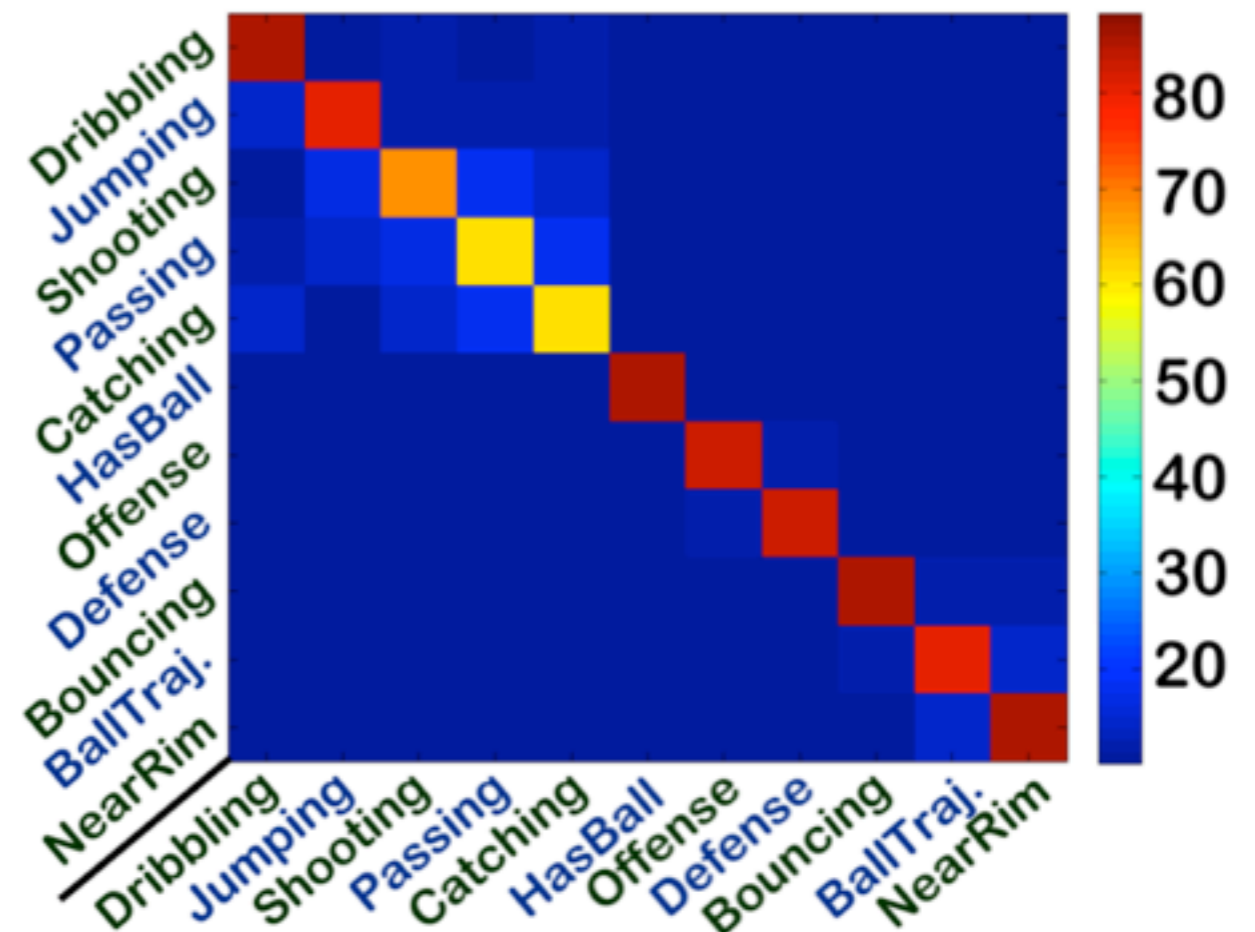


# Results -- CVPR 11

## Confusion tables



input parsing results  
of primitive actions

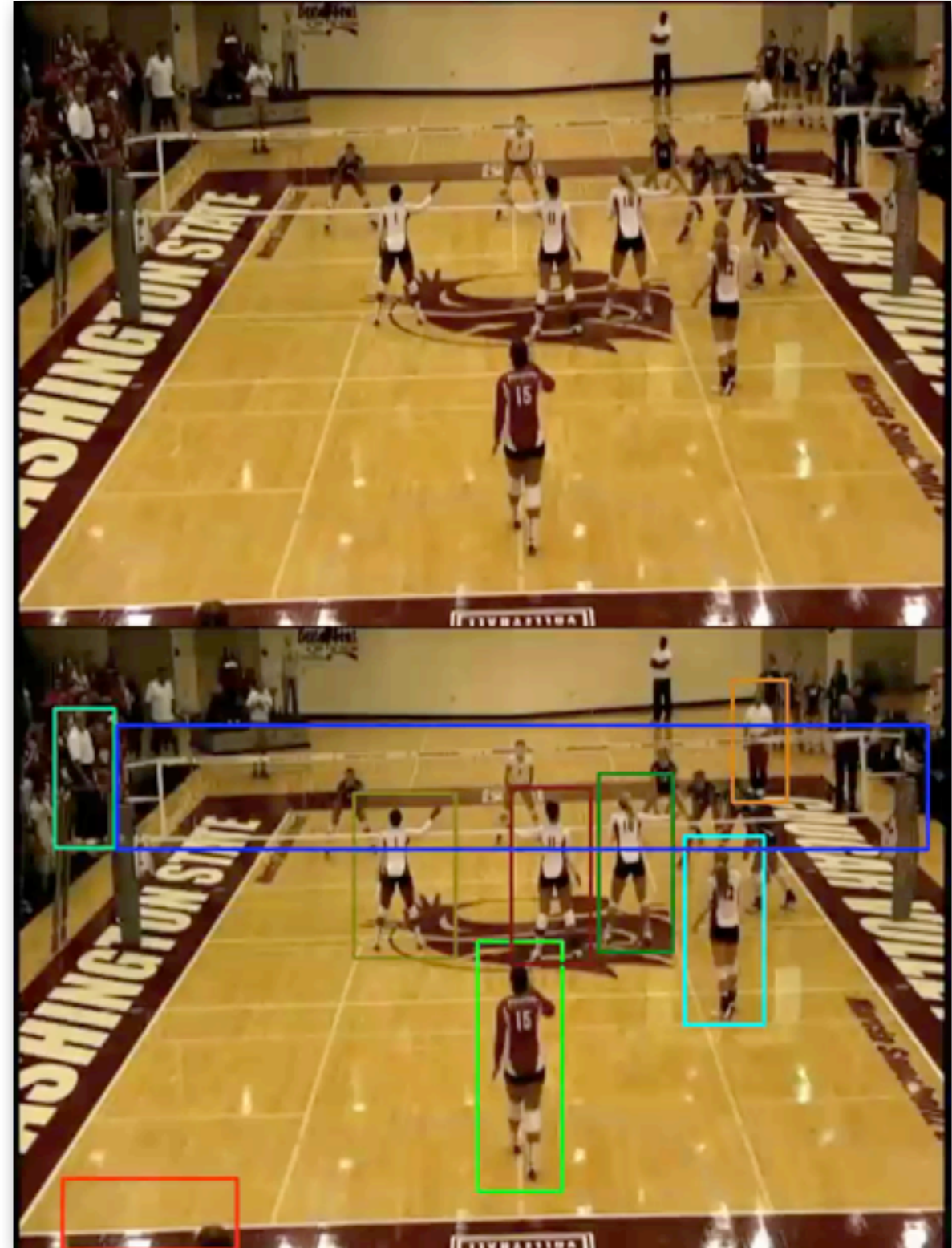


reasoning including  
higher-level events

# Summary

Reasoning helps:

- correct tracking/parsing errors
- disambiguate uncertainty
- address higher-level events









**THANK YOU**